Gems of TCS

Fine-grained Complexity

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September 14, 2022

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- For many of them, we couldn't find better algorithms in decades
- Today: Identify reason why we're stuck

Algorithmic Complexity of SAT

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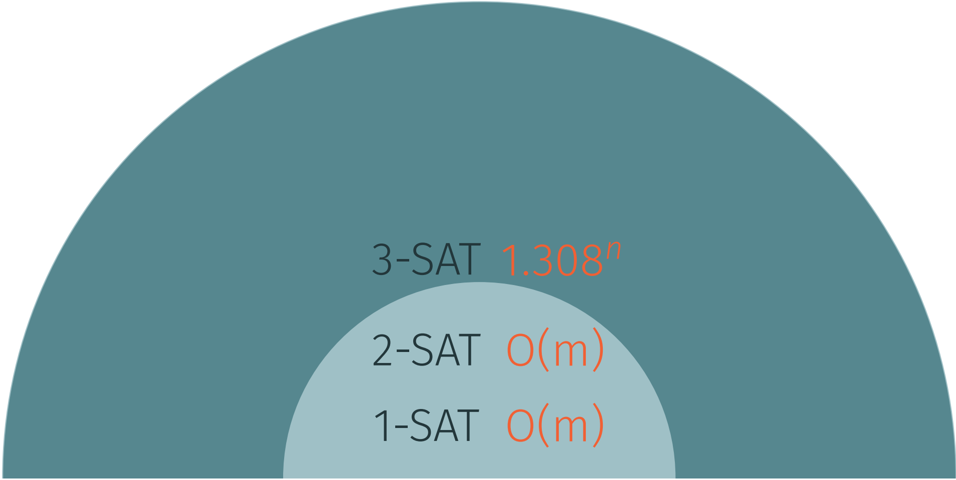

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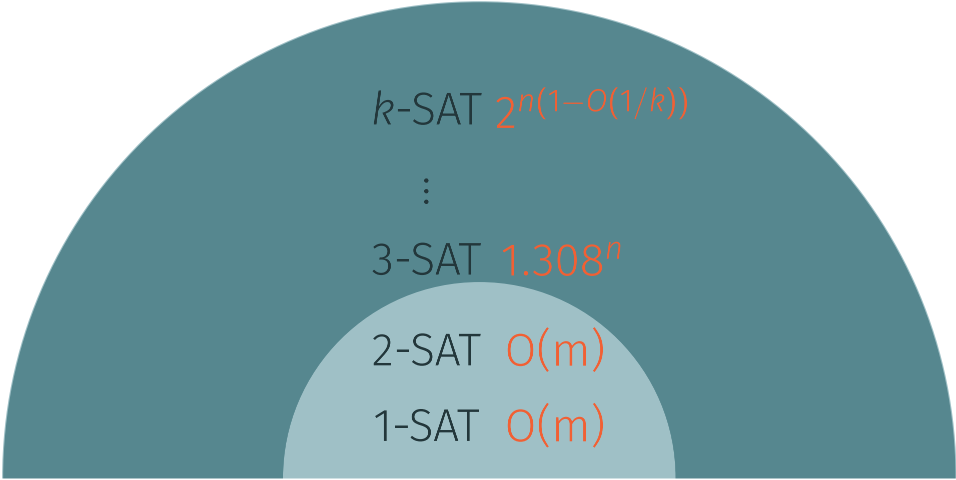

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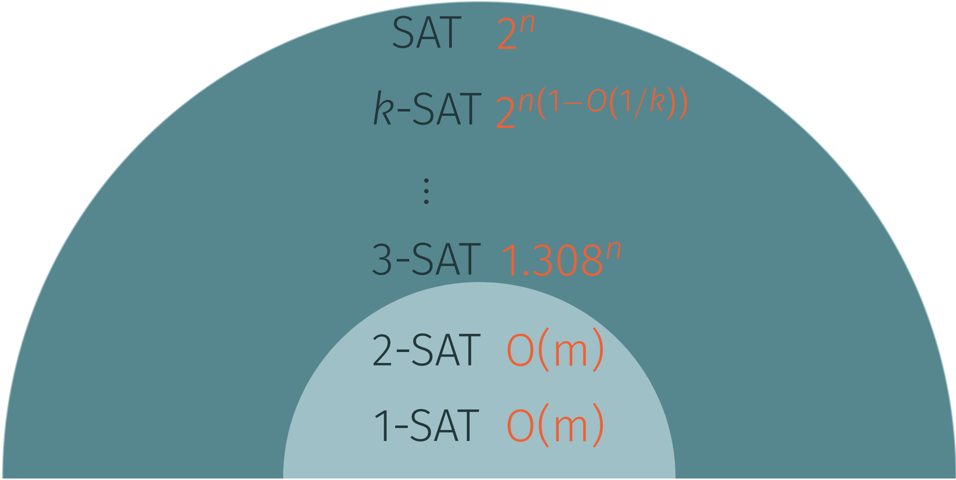

Hardness of SAT

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- Strong Exponential Time Hypothesis (SETH)

SAT requires time 2^n

Edit Distance

Edit Distance

e l e p h a n t

r e l e v a n t

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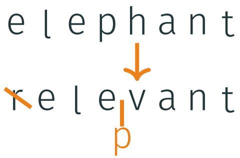

Edit Distance

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Edit Distance

$$\widetilde{O}(n^2)$$

Other Problems

Longest Common Subsequence	Orthogonal Vectors	Edit Distance
Hamming Closest Pair	All Pairs Max Flow	RNA-Folding
Regular Expression Matching	Graph Diameter	Subset Sum

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- A conjecture for each problem?
- One conjecture to rule them all?
- Fine-grained Complexity: Better-than-known algorithms for one problem would imply better-than-known algorithms for other problems

Orthogonal Vectors (OV)

Orthogonal Vectors Problem

- S, T are sets of N vectors from $\{0,1\}^d$. Are there $s \in S$ and $t \in T$ such that $s \cdot t = \sum_{i=1}^{d} s_i \cdot t_i = 0$?

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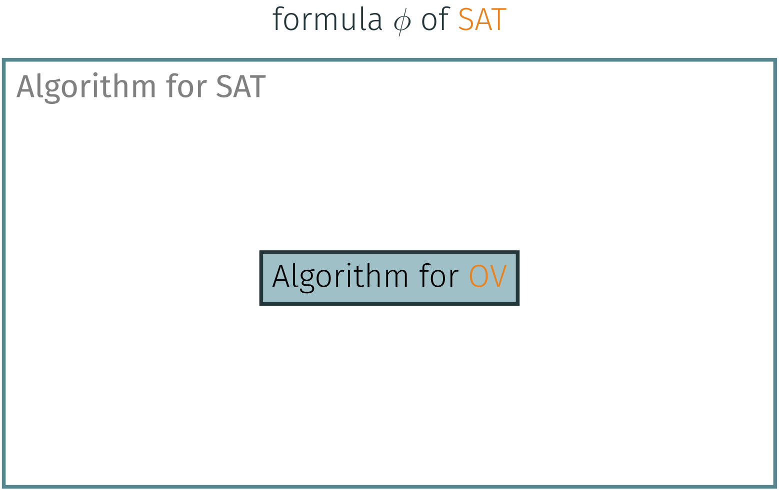
- SETH implies that OV cannot be solved in time $N^{1.99}$

Fine-grained Reductions

formula ϕ of SAT

Fine-grained Reductions



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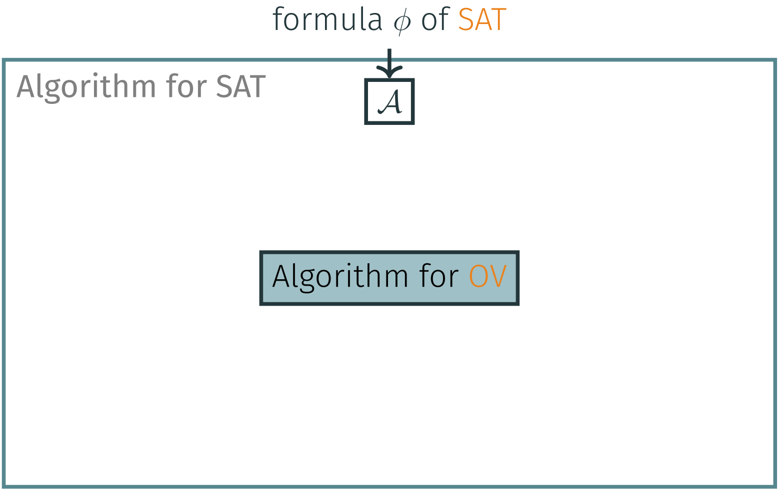

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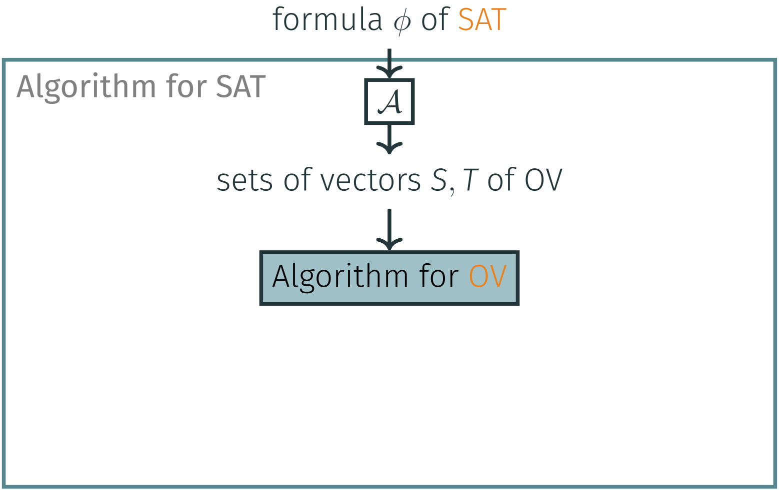

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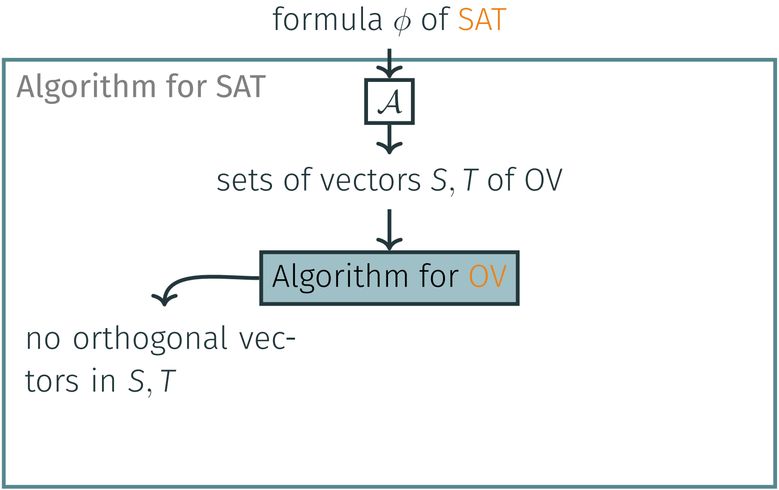

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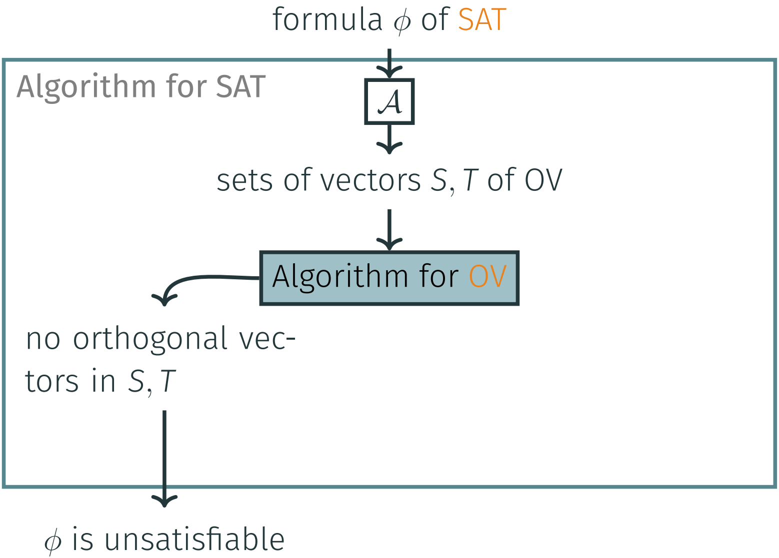

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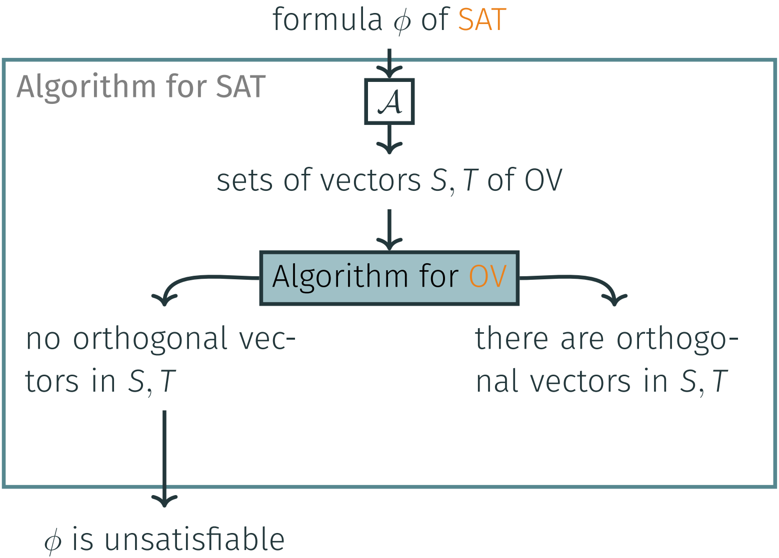

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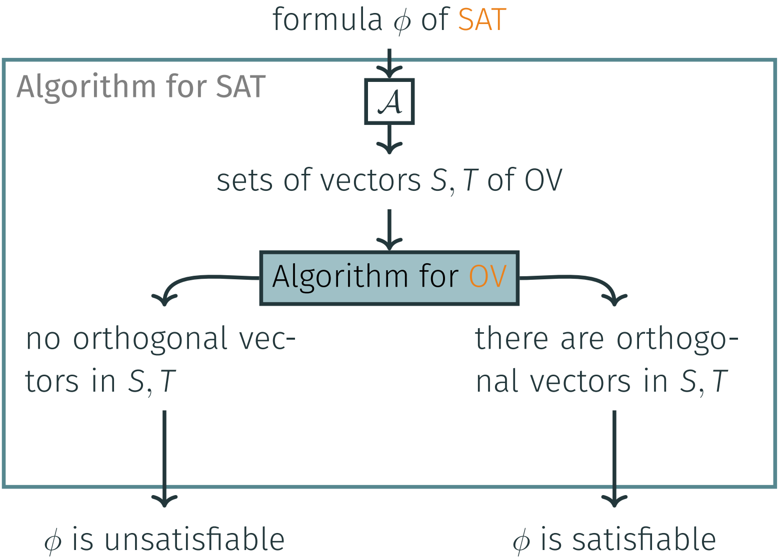

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- Given a SAT formula ϕ, split its n input variables into two sets of size $n/2$

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- $N = 2^{n/2}$

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- For an assignment $x \in \{0,1\}^{n/2}$, add $s \in \{0,1\}^m$ to S:

$$s_i = 1 \text{ iff } x \text{ doesn't satisfy clause } C_i$$

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$$N^{1.99} = (2^{n/2})^{1.99} = 2^{0.995n}$$

The Dominating Set Problem

Dominating Set

- *k-Dominating Set*: Given $G = (V, E), |V| = n$, find an $S \subseteq V, |S| = k$ such that

$$\forall v \in V \colon v \in S \text{ or } \exists u \in S \colon \ (v, u) \in E$$

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- SETH implies that k-DS cannot be solved in time $n^{k-0.01}$ for any k

SETH $\implies$ DS

Partition vars in k groups:
$\{x_1, \ldots, x_{n/k}\}, \ldots, \{x_{n-n/k+1}, \ldots, x_n\}, |\mathrm{DS}| = k$

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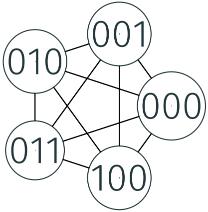
$2^{n/k}$ vertices



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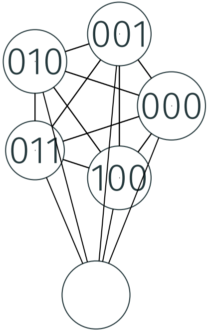

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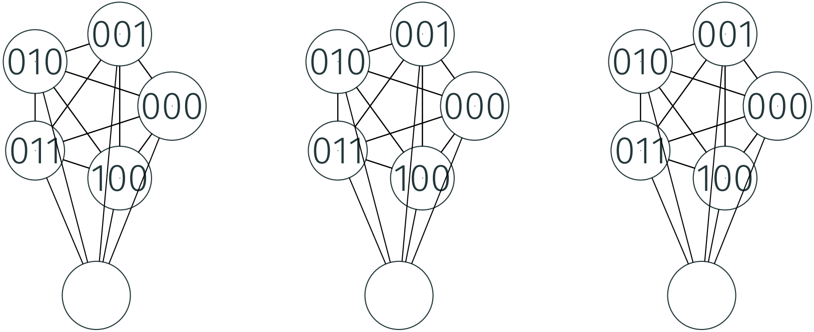

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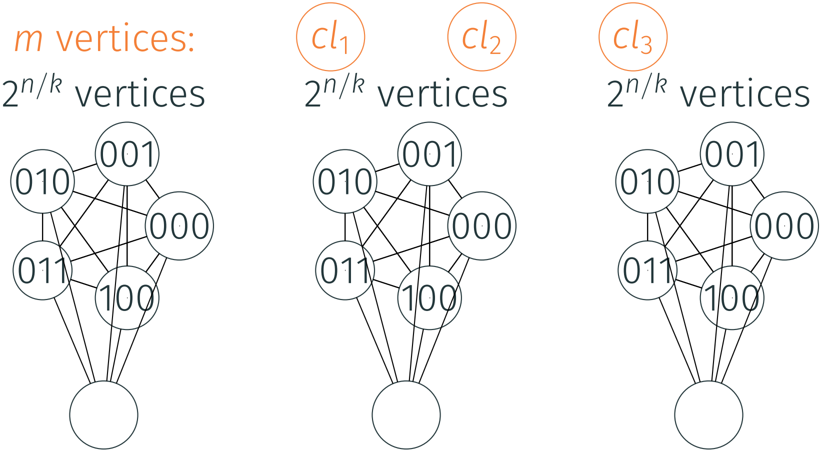

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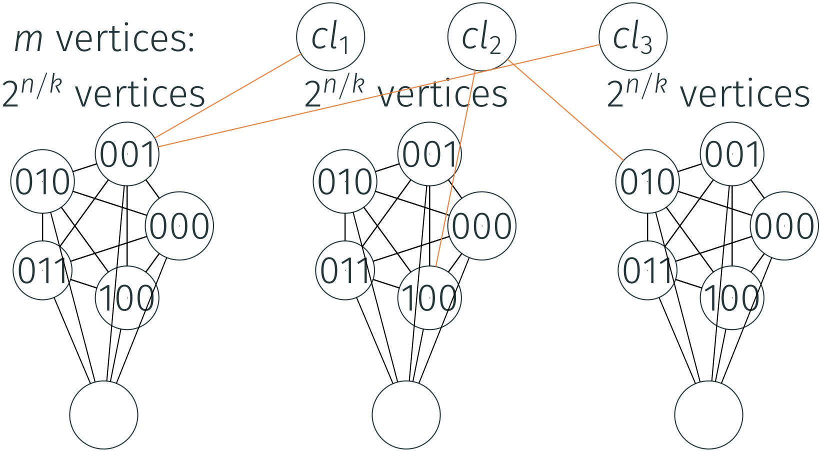

SETH $\implies$ DS

- For every k, we reduce SAT on n vertices k-DS with

$$\approx 2^{n/k}$$

vertices

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vertices
- If k-DS on N vertices can be solved in time $N^{k-0.1}$, then SAT can be solved in time
$$N^{k-0.1} = 2^{(n/k)(k-0.1)} = 2^{n-0.1n/k}$$