Gems of TCS

Streaming Algorithms

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Fruit Game

Credit: Jelani Nelson

(https://www.youtube.com/watch?v=CorP4I23wOo&t=2434s)

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- Mostly randomized algorithms

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  Sum of squares of all numbers in range $\{0, \ldots, n\}$ is $T = \frac{n(n+1)(2n+1)}{6}$

Streaming Algorithm

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  - Else count--
- Return m

Example

Another View

Misra-Gries Algorithm

Misra-Gries Algorithm

- $\mathsf{count}_1, \ldots, \mathsf{count}_k \leftarrow 0; \quad \mathsf{m}_1, \ldots, \mathsf{m}_k \leftarrow \perp$
- For each element x_i of Stream:
    - If $x_i = \mathsf{m}_j$, then count_j ++
    - Else
        - Let count_j be min in $\mathsf{count}_1, \ldots \mathsf{count}_k$
        - If $\mathsf{count}_j = 0$, then $\mathsf{m}_j = x_i; \quad \mathsf{count}_j = 1$
        - Else count_1 --, $\ldots$, count_k --
- Return $\mathsf{m}_1, \ldots, \mathsf{m}_k$

Approximate Counting

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Overview

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- Return $2^c - 1$

Analysis

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- Again, repeating Algorithm several times significantly amplifies probability of success

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- Amplify probability by Repetitions